

# The Extended Namba Problem

## §0 Introduction

Shortly after the discovery of forcing it was noticed that any regular  $\bar{\kappa} > \omega_1$  can be collapsed to  $\omega_1$  in a generic extension without adding reals. Levy showed that, if  $\kappa$  is strongly inaccessible, then every smaller cardinal can be simultaneously collapsed to  $\omega_1$ , so that  $\kappa$  becomes  $\omega_2$  in the extension. In these constructions, however, the collapsed cardinals always acquired cofinality  $\omega_1$  in the extension. Namba then showed that, assuming CH in the ground model, one can collapse  $\omega_2$  to  $\omega_1$  in such a way that it becomes  $\omega$ -cofinal in the extension. A number of interesting results were obtained by combining these collapsing techniques with iterated forcing. Fr. int. Shelah showed

that one can iterate Namba forcing out to a strongly inaccessible  $\kappa$  in such a way that  $\kappa$  becomes  $\omega_2$  and every other regular  $\tau < \kappa$  acquires cofinality  $\omega_1$ . However, the answer to the following question - which we shall here call the "extended Namba problem" remained elusive:

Assume CH. Let  $\kappa$  be strongly inaccessible. Is there a generic extension in which  $\kappa$  becomes  $\omega_2$  and in which every regular  $\tau \in (\omega_1, \kappa)$  becomes  $\omega$ -cofinal?

This seemed so difficult to us that for several years we conjectured that it was provably impossible in ZFC. Moti Gitik then refuted this conjecture by constructing a model of ZFC which admits such an extension. Gitik started with a cardinal  $\theta$  which is supercompact up to a Mahlo cardinal  $\kappa$ . He then collapsed to make  $\theta = \omega_1$

Over the resulting model he then found a generic extension in which  $\kappa = \omega_2$  and every regular  $\tau \in (\theta, \kappa)$  becomes  $\omega$ -cofinal. The question remained open, however, whether the super Namba problem always has a positive solution. We now provide an affirmative answer:

Theorem 1 Assume CH. Let  $\kappa$  be strongly inaccessible. There is a set of conditions  $\mathbb{P}$  of size  $\kappa$  s.t.  $\mathbb{P}$  is subcomplete and whenever  $G$  is  $\mathbb{P}$ -generic, then:

- $\kappa = \omega_2$  in  $V[G]$
- Every regular  $\tau \in (\omega_1, \kappa)$  becomes  $\omega$ -cofinal in  $V[G]$ .
- Every stationary subset of  $\kappa$  remains stationary in  $V[G]$

(Note By subcompleteness there are no new reals in  $V[G]$ .)

(Note  $\mathbb{P}$  will satisfy  $\kappa$ -cc, which by itself gives the last property.)

Modifying our method slightly we can also get:

Theorem 2 Let  $\kappa$  be inaccessible and assume GCH below  $\kappa$ . Let  $A_0 \subset \kappa$ .

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- $\kappa = \omega_2$  in  $V[G]$
- Let  $\tau \in (\omega_1, \kappa)$  be regular. Then  $cf(\tau) = \begin{cases} \omega_1 & \text{if } \tau \in A_0 \\ \omega & \text{if not} \end{cases}$  in  $V[G]$ .
- Every stationary subset of  $\kappa$  remains stationary in  $V[G]$  (in fact,  $\mathbb{P}$  satisfies  $\kappa$ -cc).

(Note Theorem 1 can be derived from Theorem 2, though that is not the way we do it in this paper: Using  $\omega_1$ -complete conditions we can do an Easton type collapse below  $\kappa$  s.t. GCH holds below  $\kappa$  and  $\kappa$  remains inaccessible in the extension. Since the conditions are  $\omega_1$ -complete, every regular

$\tau < \kappa$  which is collapsed in this process requires a cofinality  $> \omega_1$ .

If we then apply Thm 2 with  $A_0 = \emptyset$ , it is clear that every  $\tau \in (\omega_1, \kappa)$  which is regular in  $V$  will become  $\omega$ -cofinal.)

As one might suspect, IP is obtained as the result of an iteration - albeit an iteration in a very broad sense. The most general notion of iteration is the following:  $\langle B_i \mid i < \alpha \rangle$  is an iteration iff each  $B_i$  is a complete Boolean algebra,  $B_i \subseteq B_j$  (i.e.  $B_i$  is completely contained in  $B_j$ ) for  $i \leq j < \alpha$ , and for all limit  $\lambda < \alpha$ ,  $B_\lambda$  is generated by  $\bigcup_{i < \lambda} B_i$ . Most often  $B_\lambda$  is taken as the direct limit or the inverse limit of  $\langle B_i \mid i < \lambda \rangle$ , or as something intermediate between the two. In the present case, however, we shall often be forced to use a completion  $B_\lambda$  which has a much larger cardinality than either the direct or inverse limit. (The reasons for this are explained in §2.1.) Hence we are unable to use the iteration theorem of [SPSC]. In §1 we prove a new iteration theorem

which we shall use to ensure that each  $B_i$  in our iteration  $\langle B_i \mid i \leq \alpha \rangle$  is subcomplete. In § 2 we open with a discussion of the constraints that the iteration must satisfy. We show fr. ins. that the first stage of the iteration, whose purpose is to give  $\omega_2$  the cofinality  $\omega$ , must, in fact, make every regular  $\tau \in (\omega_1, \omega_{\omega_1})$   $\omega$ -cofinal. Hence  $B_1$  must be much larger than  $\omega_2$ . At the first limit point (and many others) there is a similar constraint on  $B_\lambda$ , which is why  $B_\lambda$  must be larger than the direct or inverse limit. In § 2.2 we then introduce the notion of "proud forcing", which will play a large role in our proof. In § 2.3 we describe the proposed iteration more precisely. The iteration is intended to prove both Theorem 1 and Theorem 2; hence an arbitrary  $A_0 \subset \kappa$  enters into its definition. We assume GCH below or only in the case  $A_0 \neq \emptyset$ .

In §3 and §4 we then fully define the iteration and verify its properties. All of the component forcings in the proof of Theorem 1 are  $\mathcal{L}$ -forcings and can be regarded as generalizations of Namba forcing. (The successor step will, in fact, be a variation on the forcing described in the appendix to [LF] §5.)

### Bibliography

[LF]  $\mathcal{L}$ -Forcing

[SPSC] Subproper and Subcomplete Forcing

(Unfortunately, I must assume a knowledge of both papers, which exist in the form of handwritten notes and can be found on my website.)