

§ 2.3 Core Mice

Def Set $\rho^\omega = \inf \{ \rho^m \mid m < \omega \}$

Let N be a mouse. N is pure iff there is a sound mouse \bar{N} s.t.

(a) N is a simple iterate of \bar{N} above $\omega_{\bar{N}}^\omega$

(b) If M, \bar{N} have a common simple iterate above $\omega_{\bar{N}}^\omega$, then M is a simple iterate \bar{N} of \bar{N} above $\omega_{\bar{N}}^\omega$.

It is clear that if N is pure and M, N have a common simple iterate above $\omega_{\bar{N}}^\omega$, then M is pure.

Note that \bar{N} is the unique sound mouse satisfying (a), since if \bar{N}' were another one, \bar{N}' would

be a simple iterate of \bar{N} above $\omega_{\bar{N}}^\omega$. But \wedge § 2.2 Lemma 1.3, as in the argument of

a proper iterate of this sort would not be sound. Hence we may define:

Def Let N be pure. The core of N ($\text{core}(N)$) is the unique sound mouse \bar{N} s.t. N is a simple iterate of \bar{N} above $\omega_{\bar{N}}^\omega$.

In this section we show that every mouse is pure. We also show that in the coiteration of two mice, at least one side must be simple. We also define a $p_N \in P_N^*$ which we call the standard parameter and show that it is preserved under simple iterations (i.e. $\pi_{NM}(p_N) = p_M$). These facts are so interwoven that we must, in effect, prove them simultaneously by a single induction on mice. The relation over which we induct is defined by:

Def $M R N$ iff M is a non simple iterate of N .

R is well founded by § 2.1 Lemma 2. We shall prove by TR -induction that N is superpure in the following sense:

Def N is superpure iff M is pure whenever N is a simple iterate of M .

We then define:

Def N is hereditarily pure (HP) iff every $M R N$ is superpure.

From now on, assume N to be HP. In the course of the induction we shall establish the above facts.

Lemma 1 Coiterate N with a mouse M . Then at least one side of the coiteration is simple.

prf.

Suppose not. Let $\langle M_i \leq \theta \rangle, \langle N_i \leq \theta \rangle$ be the coiterations. There are maximal i_M, i_N s.t. $\omega \alpha_{i_M}^M \in M_{i_M}, \omega \alpha_{i_N}^N \in N_{i_N}$.

Set $\bar{M} = M_{i_M} \upharpoonright \alpha_{i_M}^M, \bar{N} = N_{i_N} \upharpoonright \alpha_{i_N}^N$. Then $E_{\nu_{i_M}}^{\bar{M}}$ is a measure on \bar{M} where $\omega \rho_{\bar{M}}^w \leq \bar{M}$.

Thus M_θ is a simple iterate of \bar{M} above $\omega \rho_{\bar{M}}^w$. Similarly for \bar{N} . Hence:

$$(1) \bar{M} = \text{core}(M_\theta) = \text{core}(N_\theta) = \bar{N}.$$

Hence $i_N \neq i_M$, since otherwise

$$E_{\nu_i}^{N_i} = E_{\nu_i}^{M_i} \quad (i = i_N = i_M). \quad \text{Let e.g. } i_N < i_M.$$

Then $E_{\nu_{i_N}}^{N_{i_N}} = \emptyset$ and $E_{\nu_{i_N}}^{M_{i_M}} = E_{\nu_{i_N}}^{\bar{N}} \neq \emptyset$.

Contradiction!, since $\nu_{i_N} < \nu_{i_M}$.

QED (Lemma 1)

Using this we can show:

Lemma 2 If M, N have a common simple iterate, then M is HP.

proof of Lemma 2

Let $Q \leq R \leq M$. Claim Q is superpure.

Iterate M, N simply to a common M' .
Coiterate Q, M' to Q', M'' . Then:

(a) Q' is an initial segment of M''

(b) If the iteration from M' to M'' is simple, then Q' is a proper segment.

Otherwise M'' would be both a simple and non simple iterate of M' . But

then $Q' \leq R \leq N$ by (a), (b) & hence Q' is superpure. But Q' is a simple iterate of Q , since otherwise M'' would be an initial segment of Q' and both sides of the iteration would be non simple, violating Lemma 1. Hence Q is superpure. QED (Lemma 2).

Def Let M, N be mice. M is a derived iterate of N with derived iteration map π iff the following hold:

(a) M, N have a common simple iterate Q with $\text{rng}(\pi_{NQ}) \subset \text{rng}(\pi_{MQ})$

(b) $\pi = \pi_{MQ}^{-1} \pi_{NQ}$.

The derived iteration map is in fact unique, as is shown by the following fact.

Fact If M, N have a common simple iterate Q , then:

(a) If Q' is a common iterate of M, N and a simple iterate of M or N , then Q' is a simple iterate of both (hence M, N coiterate to a common Q').

(b) If $\pi = \pi_{MQ}^{-1} \pi_{NQ}$ is a derived iteration map, then $\pi = \pi_{MQ'}^{-1} \pi_{NQ}$

To see (a), coiterate Q, Q' to \tilde{Q}, \tilde{Q}' . Q, Q' are both simple iterates of a $P = M$ or N . If e.g. \tilde{Q} were a simple iterate of Q and \tilde{Q}' a non simple iterate of Q' , then \tilde{Q} , as an initial segment of \tilde{Q}' would be ~~a~~ a simple and non simple iterate of P . Contra!

To see (b) note that

$$\pi_{NQ} = \pi_{Q\tilde{Q}} \pi_{NQ} = \pi_{Q'\tilde{Q}} \pi_{NQ'}$$

and similarly for M .

Lemma 3 Let M be a derived iterate of N with map π . Then M is an iterate of N by a normal iteration and $\pi = \pi_{NM}$.

prf.

It is easily seen that N, M coiterate to a common Q and $\pi = \pi_{MQ}^{-1} \pi_{NQ}$.

Both sides of the iteration are simple and normal. Hence it suffices to prove:

Claim $Q = M$.

Suppose not, let $\langle N_{i \leq \theta} \rangle$, $\langle M_{i \leq \theta} \rangle$ be the two sides of the coiteration and let $i =$
 $=$ the least i s.t. either:

- (a) $M_i \neq M$ or;
 (b) M is not a derived iterate of N_i .

Then $i > 0$. Suppose $i = \lambda$, $\text{lim}(\lambda)$.

Then $M_\lambda = M$ since $M_j = M$ for $j < \lambda$.

But for $j < \lambda$ M is a derived iterate of N_j ; hence

$\text{rang}(\pi_{N_i Q}) \subset \text{rang}(\pi_{M Q})$ for $i < \lambda$.

Hence $\text{rang}(\pi_{N_\lambda Q}) = \bigcup_{i < \lambda} \text{rang}(\pi_{N_i Q}) \subset \text{rang}(\pi_{M Q})$. Hence M is a derived iterate of N_λ . Contr!

Thus $i = j+1$. Assume w.l.o.g. that $j = 0$. Let $\nu = \nu_0 =$ the least ν s.t. $E_\nu^N \neq E_\nu^M$.

Case 1 $E_\nu^N = \emptyset$.

Then $\pi_{N Q}(\kappa) = \kappa$ where E_ν^M is normal on κ , but $\kappa = \text{crit}(\pi_{M Q})$. Hence $\text{rang}(\pi_{N Q}) \not\subset \text{rang}(\pi_{M Q})$. Contr!

Case 2 $E_\nu^M = \emptyset$.

Then $M_1 = M$ and $\pi_{01} : N \xrightarrow{E_\nu} N_1$.

If $f \in \Gamma(N, \kappa)$, we clearly

$$\text{have } \pi_{N_1 Q}(\pi_{01}(f)(\kappa)) =$$

$$= \pi_{N Q}(f)(\kappa) = (\pi_{M Q} \pi(f))(\kappa) =$$

$$= (\pi_{M Q} \pi(f))(\pi_{M Q}(\kappa)) =$$

$$= \pi_{M Q}(\pi(f)(\kappa)). \text{ Hence}$$

$\text{rng}(\pi_{N_1 Q}) \subset \text{rng}(\pi_{M Q})$ and M is a derived iterate of N_1 . Contradiction!

Case 3 The above cases fail.

Let $X \in \mathcal{P}(K) \cap N = \mathcal{P}(K) \cap M$. Let $\pi_{NQ}(X) = \pi_{MQ}(Y)$. Then $Y \subset \pi(K) \supseteq K$ and

$$\pi_{MQ}(X) = \pi_{MQ}(Y) \cap \pi_{MQ}(K) = \pi_{NQ}(X) \cap \pi_{MQ}(K)$$

where $\pi_{MQ}(K) \supset K$. Hence:

$$\begin{aligned} X \in E_{\nu}^N &\iff K \in \pi_{NQ}(X) \\ &\iff K \in \pi_{MQ}(X) \iff X \in E_{\nu}^M. \end{aligned}$$

Thus $E_{\nu}^N = E_{\nu}^M$. Contradiction!

QED (Lemma 3)

Lemma 3.1 Let $p \in P_N^m$. Let \bar{N}, \bar{p} be s.t. $\bar{N}^m \bar{p} = N^m p$ and $\bar{p} \in P_{\bar{N}}^m$.

Then \bar{N}, N have a common simple iterate above ωp^m .
 proof.

Let $\sigma: \bar{N} \rightarrow \sum_1^{(m)} N$ s.t. $\sigma \upharpoonright \bar{N}^m \bar{p} = \text{id}$; $\sigma(\bar{p}) = p$, \bar{N} is iterable above ωp^m by the existence of σ .

Claim $E_v^{\bar{N}} = E_v^N$ for $v \leq \omega p^m$

pf. For $v < \omega p^m$ this is trivial. Now let $v = \omega p^m$. If $E_v^{\bar{N}} = E_v^N = \emptyset$ there is nothing to prove. Otherwise

κ is the largest cardinal in $\bar{N}^m \bar{p} = N^m p$, where $E_v^{\bar{N}}$ or E_v^N has critical pt. κ .

If $\omega p^m = \text{om } N$ there is nothing to prove. If not, then

$v = \kappa + N$. But then $v \in \bar{N}$, since otherwise $N^m p = \bar{N}^m \bar{p} \in N$. Hence

$v = \kappa + \bar{N}$, since $\sigma(\kappa) = \kappa$.

Hence $\sigma(v) = v$ and $\sigma(E_{\nu}^{\bar{N}}) = E_{\nu}^N$.

But $\sigma \upharpoonright R_{\bar{N}} = \text{id}$. \square ED (Claim)

But then \bar{N}, N are coiterable. Coiterate them to \bar{N}', N' . Then neither is a proper initial segment of the other, since otherwise $A \in \bar{N}$ or $A \in N$, where $A = A_N^{n,p}$ \square ED (Lemma 3.1)

Lemma 3.2 Let $p \in P_N^n$. Then $p_N^{n+m} = p_{N^m p}^m$ ($m \geq 0$). (In other words, p can be lengthened to a $q \in P_N^{n+m}$ for any $m \geq 0$).

proof.

Let \bar{N}, \bar{p} be as above & let Q be a common iterate above $w p^n$.

Then $Q^{n+\pi_{\bar{N}Q}}(\bar{p}) = \bar{N}^m \bar{p} = N^m p = Q^{n+\pi_{NQ}}(p)$

Since $\bar{p} \in R_{\bar{N}}^m$ we can lengthen it to $\bar{q} \in P_{\bar{N}}^{n+m}$. Hence $\bar{N}^{n+m}, \bar{q} =$

$Q^{n+m}, \pi_{\bar{N}Q}(\bar{q}) = Q^{n+m}, \pi_{NQ}(q) = N^{n+m}, q$

where q is the same lengthening of p . \square ED (Lemma 3.2)

Def $P_N =$ the $<_*$ -least $p \in P_N^*$,

where $<_*$ is the well ordering of $[On]^{<\omega}$ defined by:

$$a <_* b \iff \forall \nu (a \setminus (\nu+1) = b \setminus (\nu+1) \wedge \nu \in b \setminus a).$$

If we set: $P_N^i = P_N \cap [\omega p_N^{i+1}, \omega p_N^i)$ ($i < \omega$), then it follows from Lemma 3.2 that:

$$P^i = \text{the } <_*\text{-least } p \in P_N^i \cup P^i.$$

→

We call p_N the standard parameter of N . The main lemma in this section states the preservation property of a standard parameter: $\pi_{NM}(p_N) = p_M$.

In preparation for this we define the Friedman witnesses: *

* / This is something of a misnomer, since these witnesses, though discovered independently by Friedman, turn out to have been defined earlier by Mitchell.

An the following we adopt the convention that $P^i = P_N \cap [\omega p_N^{i+1}, \omega p_N^i)$ where $P \subset On_N$ is finite, and $P^i = \langle P^0, \dots, P^{i-1} \rangle$.

Def Let $\nu \in p^i = p_N^i$ ($i < \omega$).

The Friedman witness for ν ($W = W_N^\nu$) is defined as follows: Let

$$X = h_{N^i, P^i}(\nu \cup (p^i \setminus (\nu+1)))$$

and $\bar{W} \cong (N^i, P^i \upharpoonright X)$ where \bar{W} is transitive. W, \bar{p} are the unique objects s.t. $W^i P = \bar{W}$ and $\bar{p} \in R_W^i$.

(Note that $\bar{p} = p_W^i$, and $\omega p_W^{i+1} \leq \nu$).

The following lemma marks our first use of (c) in the def. of premouse.

Lemma 4.1 Let $\nu \in p = p_N$. Then $W_N^\nu \in N$.

proof

Let $\bar{\sigma}: \bar{W} \xrightarrow{\sim} (N^i, P^i \upharpoonright X)$, $\sigma \supset \bar{\sigma}$,

with $\sigma: W \rightarrow \sum_1^{\omega} N$, $\sigma(\bar{p}) = p = p_N^i$

(where \bar{w}, W, \bar{p} are as in the above def.).

The existence of σ tells us that

(1) W is iterable above $\omega p_W^{i+1} \leq \nu$.

(2) W, N are coiterable.

proof. of (2)

Clearly $E_{\xi}^N = E_{\xi}^W$ for $\xi < v$. Coiteration, could only fail if $E_v^N \neq E_v^W \neq \emptyset$.

Assume $E_v^W \neq \emptyset$. If $v \in W$, then

$$E_v^W = E_{\sigma(v)}^N \cap J_v^{E^N} \text{ where } E_{\sigma(v)}^N \text{ is a}$$

measure on the same κ as E_v^W . Hence,

since W is a pm, the def. of pm tells us that $E_v^N = E_{\sigma(v)}^N \cap J_v^{E^N} = E_v^W$.

If $v \notin W$, then $E_v^W = E_{\alpha}^N \cap J_v^{E^N}$ where

$W \setminus \alpha = \emptyset \cap N$ and we apply the same

reasoning. QED (2)

Note that by the proof of (2), the W side of the coiteration is above v , and the N side is above ω_p^{i+1} .

Coiterate N, W to N', W' .

(3) $N' \not\subseteq W'$. Then N' is a simple iterate of N

Suppose not. Imitating the W side of the coiteration, we

get an iterate N'' of N' above v

and a map $\sigma' : W' \xrightarrow{\sum_1^{\omega} \pi} N''$

$$\text{s.t. } \sigma' \pi_{W, W'} = \pi_{N, N''} \sigma$$

$$\begin{array}{ccc} W & \longrightarrow & W' \\ \sigma \downarrow & \nearrow N' & \downarrow \sigma' \\ N & \longrightarrow & N'' \end{array}$$

Hence $\sigma' \pi_{NN'} : N \xrightarrow{\Sigma^*} \sigma'(N')$ where $\sigma'(N')$ is a proper initial segment of N'' , hence a non simple iterate of N . Contr! by § 2.1 Lemma 3. QED (3)

(4) $N' \neq W'$ if N' is a simple iterate of N .

pf.

The iteration from W to W' is simple, since otherwise N'' is a non simple iterate of N and we get the same contradiction as in (3). But then $p_N^{i+1} = p_{N'}^{i+1} = p_{W'}^{i+1} = p_W^{i+1}$, since both iterations are above ωp^{i+1} . Let $A \subset \omega p^{i+1}$ be $\Sigma_1(\bar{w})$ s.t. $A \notin N$.

Set: $\bar{p}^* = \bar{\sigma}^{-1}(p^*)$ where $p^* = p_N^i \setminus (v+1)$

Since $h_{\bar{w}}(v \cup \bar{p}^*) = \bar{w}$, there must be $\vec{s} < v$ s.t. A is $\Sigma_1(\bar{w})$ in $\{\vec{s}\} \cup \bar{p}^*$

Hence A is $\Sigma_1(N^i, P^i)$ in $\{\vec{s}\} \cup P^*$, contradicting the def of P_N .

QED (4)

Hence:

(5) Either $W' \in N'$ or $W' = N'$ and the iteration from N to N' is non simple.

We now prove $W \in N$ by cases:

Care 1 There is $\beta \in N$ s.t. $\bar{\beta}^N > \nu$.

Set: $D = \{ \langle \bar{z}, \bar{s} \rangle \in V^2 \mid h_{\bar{w}}(\bar{z}, \bar{p}^*) \in h_{\bar{w}}(\bar{s}, \bar{p}^*) \}$.

Clearly, ν is p.m. closed & hence W can be constructed from D inside any admissible set. It suffices to prove:

Claim $D \in N$,

since then $D \in J_{\beta}^{EN}$ where J_{β}^{EN} is admissible p.f. of Claim.

If $W' \in D'$, then $D \in N'$ since D is $\Sigma^*(W')$. Hence $D \in N_1$ since N' is an iterate of N' above ν ($\langle N_{\beta} \mid \beta \leq \theta \rangle$ being the normal iteration. But

Then $D \in N$, since N_1 is either an it. of N above ν or the Σ^* ultrapower of an NId by E_{ν}^N (hence $\mathcal{P}(\nu) \cap N_1 \subset N$).

Now let $W' = N'$. Then N' is a non simple iterate of N and there is a maximal $\beta < \theta$ s.t.

$$\pi_{\beta} : (N \mid \beta) \xrightarrow{E_{\beta}} N_{\beta+1} \text{ with } w_{\alpha} \in N_{\beta}.$$

But N' is a simple iterate of $N_{\zeta+1}$ above ν . Hence $D \in \underline{\Sigma}^*(N_{\zeta+1})$ by §2.2 Lemma 5. But either $\kappa_{\zeta} \geq \nu$ or $\zeta = 0$ and $\nu = \kappa + N$. In the former case, $D \in \underline{\Sigma}^*(N_{\zeta} \text{ id}) \subset N_{\zeta}$ and it follows as before that $D \in N$. In the latter case we note that there is $\langle f_{\zeta} \mid \zeta < \nu \rangle \in N$ w.t. $f_{\zeta} : \kappa \rightarrow \nu$, $f_{\zeta} \in N_{\zeta}$ and $\pi(f_{\zeta} \mid \kappa) = \zeta$ for $\zeta < \nu$. By §1.3 Cor 2.5.1 it follows easily that $D \in N$. QED (Case 1)

Case 2 Case 1 fails and $\nu \in W$. Then $\sigma(\nu) = \nu$, since otherwise ν is a cardinal in W and $\sigma(\nu) > \nu$ would be a cardinal in N . But then $i = 0$ and $W = N \cap X = \langle J_{\delta}^{EN}, E_{\alpha}^N \cap J_{\delta}^{EN} \rangle$ for some δ , where $\alpha \delta = \text{On} \cap N$. But then $W = W'$, $N = N'$ and $W \in N$ by (5). QED (Case 2)

Case 3 $\omega \nu = \text{On} \cap W$
Then $W = \langle J_{\nu}^E, E_{\alpha}^N \cap J_{\nu}^E \rangle \in N$.

QED (Lemma 4.1)

Lemma 4.1 makes it clear in what sense $W = W_N^v$ is a "witness" to $v \in P_N$.

If $v \notin P_N$, then there would be an $A \in \omega p^{i+1}$ s.t. $A \notin N$ which is $\Sigma_1(N^{i+1}, P_N)$ in parameters from $v \cup P_N \setminus (v+1)$. Hence $A \in \Sigma^*(W) \in N$. Contradiction!

Now suppose $\pi: N \xrightarrow[E_3]{*} M$. Let $\tilde{P} = \pi(P_N)$. Then $\tilde{P} \in P_M^*$. If we can show that each $v \in \tilde{P}$ is "witnessed" inside M in the same way, we can conclude that $\tilde{P} = P_M$. Following this strategy we prove:

Lemma 4.2 Let $\pi: N \xrightarrow[E_3]{*} M$.

Then $\pi(P_N) = P_M$.

proof of Lemma 4.2.

Let $p = p_N$, $q = \pi(p)$. Then $q \in P_M^*$.

Let $\tau \in q$. In particular, suppose $\tau \in q^i = q \cap [\omega p_M^{i+1}, \omega p_M^i)$. Then

$\tau = \pi(\nu)$ where $\nu \in p^i$. Set

$$\bar{N} = N^i p^i, \quad \bar{M} = M^i q^i,$$

$$p^* = p^i \setminus (\nu + 1), \quad q^* = q^i \setminus (\tau + 1).$$

By our above remarks it suffices to show that there are $Q, \bar{q}^* \in M$ s.t.

$$(1) \wedge \vec{s} < \tau (Q \models \varphi(\vec{s}, \bar{q}^*) \leftrightarrow \bar{M} \models \varphi(\vec{s}, q^*))$$

for Σ_1 formulae φ , since there are Q witness the minimality of q in P_M^* . Let E_3 be a measure on κ .

Case 1 $\kappa < \omega p_N^{i+1}$

Then $\pi \upharpoonright \bar{N} : \bar{N} \xrightarrow{\Sigma_2} \bar{M}$. Since

$$(2) \wedge \vec{s} < \nu (\bar{W} \models \varphi(\vec{s}, \bar{p}^*) \leftrightarrow \bar{N} \models \varphi(\vec{s}, p^*)),$$

the claim holds with $Q = \pi(\bar{W})$.

Case 2 $\omega p_N^{i+1} \leq \kappa < \omega p_N^i$

Then $\pi \upharpoonright \bar{N} : \bar{N} \xrightarrow{\Sigma_1} \bar{M}$ cofinally.

Case 2.1 There is $X \in \bar{N}$ s.t. $\text{rng}(\bar{\sigma}) \subset X$
(where $\bar{\sigma} : \bar{W} \rightarrow \bar{N}$ is defined as before).

Take X as transitive.

$$(3) \Lambda_{\vec{z}}^{\vec{z}} < \nu (\bar{N} \models \varphi(\vec{z}, p^*) \rightarrow \bar{W} \models \varphi(\vec{z}, \bar{p}^*))$$

for Σ_1 formulae φ . (3) is $\Pi_1(\bar{N})$ in the parameters \bar{W}, \bar{p}^*, p^* . Moreover:

$$(4) \Lambda_{\vec{z}}^{\vec{z}} < \nu (\bar{W} \models \varphi(\vec{z}, \bar{p}^*) \rightarrow (\bar{N} \upharpoonright X) \models \varphi(\vec{z}, p^*))$$

which is $\Pi_1(\bar{N})$ in $\bar{W}, (\bar{N} \upharpoonright X), \bar{p}^*, p^*$. It follows easily that the claim holds for $Q = \pi(\bar{W})$.

Case 2.2 Case 2.1 fails

Let $\Gamma = \{u \in W \mid u \text{ is transitive}\}$.

Then $\bar{N} = \bigcup_{u \in \Gamma} \bar{\sigma}(u)$. For $u \in \Gamma$:

$$(5) \Lambda_{\vec{z}}^{\vec{z}} < \nu ((\bar{W} \upharpoonright u) \models \varphi(\vec{z}, \bar{p}^*) \leftrightarrow (\bar{N} \upharpoonright \bar{\sigma}(u)) \models \varphi(\vec{z}, p^*))$$

which is $\Pi_1(\bar{N})$ in $\bar{W} \upharpoonright u, \bar{N} \upharpoonright \bar{\sigma}(u), \bar{p}^*, p^*$.

But $\bar{M} = \bigcup_{u \in \Gamma} \pi \bar{\sigma}(u)$. Hence (1)

holds with $Q = \pi(\bar{W}) \upharpoonright \bigcup_{u \in \Gamma} \pi(u)$.

But $Q \in M$, since $Q = \langle J_{\bar{r}}^A, B \cap J_{\bar{r}}^A \rangle$

where $\pi(\bar{W}) = \langle J_{\bar{r}}^A, B \rangle$ and

$$\bar{r} = \sup \pi'' \text{On} \cap \bar{W}.$$

Case 3 $\omega p_N^i \leq \kappa$.

Then the claim holds with $Q = \bar{W}$.

QED (Lemma 4.2).

It follows easily from the proof of Lemma 4.2 that:

Cor. 4.3 Let $\pi: N \xrightarrow[E_3^*]{} M$, $v \in P_N$.

Then either $W_M^{\pi(v)} = \pi(W_N^v)$ or

else $(0_N \cap W_M^{\pi(v)}) \in \pi(W_N^v)$.

We use these lemmas to prove:

Lemma 5 Let M be a simple iterate of N

Then $\pi_{NM}(P_N) = P_M$.

Proof.

Let $\langle N_i \mid i \leq \theta \rangle$ be the iteration

from N to M with int. maps π_{ij} .

By induction on i we prove:

Claim $\pi_{hi}(P_{N_h}) = P_{N_i}$ for $h < i$.

Only the limit case is problematic.
 Let the claim hold for $i < \lambda$.

By Cor 4.3 there is an $i_0 < \lambda$
 s.t. $\pi_{ij} (W_{N_i}^\nu) = W_{N_j}^{\pi_{ij}(\nu)}$ for $\nu \in P_{N_i}$

when $i_0 \leq i \leq j < \lambda$. Set

$P = \pi_{i\lambda} (P_{N_i})$ ($i < \lambda$). Set

$W^\nu = \pi_{i\lambda} (W_{N_i}^{\pi_{i\lambda}^{-1}(\nu)})$ ($i_0 \leq i < \lambda$)

for $\nu \in P$. It follows easily
 that the W^ν witness the

$<_*$ -minimality of p in $P_{N_\lambda}^*$.

QED (Lemma 5)

We are now ready to prove:

Lemma 6 N is pure. Moreover if $\rho_N^\omega = \rho_N^m$ and \bar{N}, \bar{p} are defined

by: $\bar{N}^m \bar{p} = N^m p^m$, $\bar{p} \in \mathbb{R}_{\bar{N}}^m$,

where $p = p_N$, then $\bar{N} = \text{core}(N)$,

proof.

Let $\sigma: \bar{N} \rightarrow \sum_1^{(m)} N$ s.t. $\sigma \uparrow \bar{N}^m \bar{p} = \text{id}$ and

$\sigma(\bar{p}) = p^m = \langle p^0, \dots, p^{m-1} \rangle$. Then

$\sigma: \bar{N} \rightarrow \sum_* N$. By Lemma 3.1 \bar{N} and N

have a common simple iterate Q

above $\omega \rho_N^m = \omega \rho_{\bar{N}}^m$. Hence \bar{N} is HP.

Clearly $p_{\bar{N}} = \bigcup_{i < m} \bar{p}^i$ and

$\pi_{\bar{N}Q}(p_{\bar{N}}) = \pi_{NQ}(p_N) = p_Q$. Since

$p_{\bar{N}} \in \mathbb{R}_{\bar{N}}^*$ and $\pi_{\bar{N}Q} \uparrow \omega \rho_{\bar{N}}^m = \pi_{NQ} \uparrow \omega \rho_N^m$,

we conclude: $\sigma = \pi_{NQ}^{-1} \pi_{\bar{N}Q}$.

Hence by Lemma 3, N is an iterate of \bar{N} and $\sigma = \pi_{\bar{N}N}$. Now suppose

that M, \bar{N} have a common simple iterate above $\omega \rho_{\bar{N}}^m$.

The same analysis then shows that M is a simple iterate of \bar{N} .
 It remains only to show:

Claim \bar{N} is round.

prf.

It suffices to show for $i < n$ that $P_{\bar{N}^i, \bar{P}^i} = R_{\bar{N}^i, \bar{P}^i}$. Suppose not, let i be the least counterexample.

Let q be $<_*$ -minimal in $P_N \setminus R_N$, where $N = \bar{N}^i \cap \bar{P}^i$.

Then $p <_* q$, where $p = \bar{P}^i$.

Let N', q' be s.t. $N'q' = Nq$ and $q' \in R_{N'}$. Let

$\sigma: N' \rightarrow_{\Sigma_1} N$ be def. by

$\sigma \upharpoonright N'q' = \text{id}$, $\sigma(q') = q$. Then

(1) $N = \bigvee_{p <_* q} q = h(i, \langle \vec{3}, p \rangle)$

for some $\vec{3} < \omega_f_N$, since

$p \in R_N$. Hence there is

$p' \in N'$ s.t. $p' <_* q'$ and
 $q' = h_{N'}(i, \langle \vec{s}, p' \rangle)$. Then,
~~rather~~ $\sigma(p') \in P_N$ and
 $\sigma(p') <_* q$. But $\sigma(p') \notin \Pi_N$,
 since $\sigma \neq \text{id}$. Contradiction!
 by the minimality of q .

QED (Lemma 6)

But if N is a simple iterate of M ,
 then M is HP & hence pure by
 Lemma 6. Thus... N is superpure.
 This completes the induction,
 showing that all mice are pure.
 Hence the previous lemmas hold
 for all mice.